Self-Protection Strategies

Tunneling, Armored, and Retro Viruses

CS4400/7440

Anti-anti-virus Techniques

- Virus writers have devised numerous methods of resisting anti-virus software and making life difficult for anti-virus researchers
- We will examine four categories of virus self-protection in coming weeks:
 - tunneling,
 - armor,
 - retroviruses, and
 - encrypted viruses of several types
- Reading Assignment: Chapter 6 of Szor.

Tunneling Viruses

Recall the DOS IVT (interrupt vector table) and the technique of interrupt hooking:



Background: Chaining Interrupt Handlers

- Interrupts contain address pointing to interrupt vector
 - Interrupt vector contains addresses of interrupt handlers.
- If more devices than elements in interrupt vector, then chain:
 - List of handlers for given address traversed to determine the appropriate one.

vector number	description
0	divide error
1	debug exception
2	null interrupt
3	breakpoint
4	INTO-detected overflow
5	bound range exception
6	invalid opcode
7	device not available
8	double fault
9	coprocessor segment overrun (reserved)
10	invalid task state segment
11	segment not present
12	stack fault
13	general protection
14	page fault
15	(Intel reserved, do not use)
16	floating-point error
17	alignment check
18	machine check
19–31	(Intel reserved, do not use)
32–255	maskable interrupts

Pentium Processor Event-Vector Table

Hooking an Interrupt

- 1. Get location/length of IDT using Intel **sidt** instr.
 - SIDT (Store Interrupt Descriptor Table) stores contents IDTR (Interrupt Descriptor Table Register) register,
 - which is a selector that points into the Interrupt Descriptor Table.
- 2. Each descriptor is 8 bytes: Index into the Table by 8n bytes to change interrupt n
- This descriptor contains the address of the Ring0 code to run for interrupt n
- This address is changed to point to hooking code
 - Additional work to chain

Interrupt Hooking

Interrupt hooking IS a legitimate technique,

• e.g. a disk compression utility might need to intercept disk accesses to compress and decompress on the fly:



Anti-virus Interrupt Monitors

When an anti-virus program executes at boot-up time, it installs a monitor that lengthens the call chain even more:



 The AV monitor checks to see if it is first on the call chain.

 If so, calls the saved address for the next item on the chain (in this case, the compression handler).

Detecting the Interrupt Hooking Virus

However, if a virus has hooked the interrupt, then the anti-virus monitor code detects that it is not being called directly from the IVT:



The AV monitor now begins virus disinfection.

Tunneling Viruses

A tunneling virus defeats the anti-virus monitor by following the interrupt call chain until it finds the end, installing itself there instead of at the beginning:





 The AV monitor now finds itself pointed to directly from the IVT and finds nothing to disinfect.

Tunneling Methods

- The process of following the interrupt call chain is called *tunneling*, because the virus is trying to locate itself in the system in a place that is beneath the vision of the anti-virus software
- How can a virus follow the call chain?
 - Emulation (sophisticated and costly)
 - Stepping through instructions in debug mode
 - In DOS, scanning all of memory to find the code that calls the BIOS handler, which must be the end of the chain

Defeating Tunneling Viruses

- The AV monitor
 - can scan in both directions and record the call chain for later checking
 - scan for virus code patterns throughout all the handlers in the call chain,
 - in case the virus had already tunneled down the chain before the AV software was installed
 - removes the virus handler when it is detected

Interrupt Wars

- An interrupt hooking virus usually has a memory-resident file infector component in addition to the interrupt handler; the handler calls the infector
- The memory-resident component can detect that the handler has been removed, and can re-install it at the end of the call chain
- The AV monitor will detect the new virus handler and remove it again; this *interrupt war*, carried on while interrupts are being processed, can make a system unstable
- Solution: find and remove the memory-resident code immediately before removing the handler

Armored Viruses

- An armored virus makes it difficult for anti-virus professionals to detect and analyze its functions
- Anti-virus professionals use a variety of detection and analysis tools:
 - Disassemblers
 - Debuggers
 - Emulators
 - Heuristic analyzers
 - Goat files
- Armored viruses try to make each of these tools ineffective or more difficult to use

Armored Viruses

- Armored virus techniques fall naturally into five categories, corresponding to the five tools they are designed to combat:
 - Anti-disassembly
 - Anti-debugging
 - Anti-emulation
 - Anti-heuristics
 - Anti-goat

Anti-Disassembly

- The broadest category of techniques that make disassembly difficult are the virus code encryption techniques, which we will study separately for several weeks starting next week. Other techniques:
 - Encrypted data
 - Code obfuscation
 - Using checksums
 - Compressed code
- We will examine each of these briefly

- The virus encrypts its data and decrypts it as it is used
- The encryption and decryption code is clearly visible, so it is straightforward to figure out
- BUT, when viewing the code in a disassembler, the data is garbled
- Labor-intensive: The anti-virus software engineer is <u>slowed down</u> by the need to emulate code, write a decryption utility program and paste data into it, etc.

Encrypted Data Example

- The Fix2001 worm attacked Windows 95 systems in 2001
- The worm sent stolen accounts and passwords by email back to a free email address (e.g. hotmail.com) obtained with a false identity
- The worm author did not want the email address to be readable to a disassembler
- The address was in a constant data section that was encrypted
- Stepping through a debugger to watch the data be decrypted slows down the analysis

Code Obfuscation

- We saw a DOS example two weeks ago that used a jump into the middle of a previous instruction
- Some obfuscation merely injects no-ops, do-nothings (e.g. add eax, 0)
 - Regular expression matching can filter these out
 - Analysis is not slowed much by these instructions
- It is slower to analyze code with roundabout computations, computed jump addresses rather than direct jumps, etc.

Obfuscated Computation

- Example from Szor text, p. 223:
- > Straightforward code to write 256 bytes into a file:

mov	CX,	100h	;	100h = 256 bytes to write
mov	ah,	40h	;	40h = DOS function number
int	21h		;	Invoke DOS handler

Convoluted code to do the same thing:

```
mov cx,003Fh ; cx = 003fh
inc cx ; cx = 0040h
xchg ch, cl ; swap ch, cl (cx = 4000h)
xchg ax, cx ; swap ax, cx (ax = 4000h)
mov cx, 0100h ; cx = 100h
int 21h ; Invoke DOS handler
```

Anti-Disassembly Checksums

- Straightforward code to match an imported function prototype, from the exported functions list in DLL,
 - to decide which system functions to infect,
 - might loop through the DLL function names list and
 - compare each function name to a constant string, e.g. (in C pseudocode),

for (each prototype in DLL export table)

if (0 == strcmp(name, "GetFileHandle(int)"))

infect(current export table address);

endfor

• Easy to read in the disassembled code;

good disassembler can even search and find the string "GetFileHandle" if the antivirus researcher already suspects that is the function being infected

Checksums cont'd.

Instead, the virus could compute a checksum over the ASCII bytes of the two strings, store one as a constant, and compare the checksums for equality:

```
int ConstantName = 0x89f7e5b2; /* Computed by virus writer */
for (each prototype in DLL export table)
    int foo = checksum(name);
    if (foo == ConstantName)
        infect(current export table address);
endfor
```

- > This code no longer reveals the API name to a reader
- Labor Intensive: Anti-virus researcher must now <u>step</u> through the checksum computation to figure out what is going on
 - i.e., impedes the analysis
- Similar idea to encrypting data

Anti-Disassembly Compression

- A virus can be stored using a compression algorithm, and decompressed during execution by a decompression code at the beginning of the virus
- As with encrypted data, the compression algorithm is exposed, but examination of disassembled code is greatly slowed down
- Anti-virus researcher might need to emulate the code, or step through it in a debugger

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Anti-Debugging

- We have seen that anti-disassembly techniques might drive an anti-virus researcher to step through virus code in a debugger
- The next step in the escalating war between the virus and anti-virus communities is the development of virus code that resists being executed in a debugger

Interrupts I and 3 are used often in x86 debugging

- **INT** 1 places the CPU in single-step mode
- INT 3 is inserted into the code by the debugger to set a breakpoint
- First anti-debugging technique: Hook these two interrupts
 - Anti-virus code must be used to unhook the interrupts before debugging can proceed
- Next anti-debugging technique: use a checksum to defeat INT 3 breakpoints

If the virus code computes a checksum over a critical range of its code, stores that checksum in a constant, and then recomputes the checksum as it runs, it can detect the change when the INT 3 instruction is injected into that code range, and just abort

The virus now runs successfully without a debugger, but aborts when breakpoints are inserted

 Note that a code emulator, as opposed to a debugger, does not insert breakpoints and is not defeated by this technique

- Next technique: Detecting changes in the state of the stack during single-step mode
- A single-step debugger places a state record on the stack, and updates it after each step, to record the current IP (instruction pointer) value and the contents of the FLAGS register
- The virus code can examine the stack and see changing values where they would not change during normal (non-debugger) execution, and abort

Detecting Stack Changes

- In single-step debug mode, the debugger, after each instruction, saves state change information in a record that is placed just beyond the top of the stack.
- This location will be trampled by any user program instructions that change the stack, but the old info is not needed after an instruction is executed.



Detecting Stack Changes

Without single-step debug state changes, a location on the stack will remain unchanged until an instruction changes it, but will be changed by the debugger after every instruction during single-step debug:

mov bp,sp ; bp gets current stack pointer
push ax

pop ax ; old pushed value still at [bp-2]
; which is beyond current stack
cmp word ptr [bp-2],ax ; equal if no debugger
jne DEBUG ; debugger detected! Go abort!

If the virus encrypts part of its code, it can use the stack pointer in the decryption routine

• When the debugger uses the stack to save state, it will change the stack pointer and cause the decryption to fail

Virus now executes only without a debugger

- Disabling the keyboard interrupt during virus execution prevents the AV researcher from using a debugger
- Other techniques are listed in Szor, 6.2.7

Detecting Debuggers

 On Win32 operating systems, an API is available: IsDebuggerPresent()

Virus can just abort if TRUE

- Many debuggers set registry keys when they are active
- Viruses can scan memory for debugger code

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Anti-Emulation Techniques

▶ em u la tor **noun**

- I : one that emulates; imitator
- 2: hardware or software that permits programs written for one computer to be run on another computer
- Emulators <u>approximate</u> the behavior of their target
 - Why is this?
 - Emulation runs much more slowly than actual machine

Anti-Emulation Armoring:

- Floating point code
- Take advantage of the approximation be exceptional
- Time/logic bombs
- Dynamic code length

Anti-Emulation: Floating Point Code

- If armored viruses have driven AV researchers away from disassemblers and debuggers,
 - next step for virus writers is to impair use of code emulators
- Early emulators only kept track of the integer CPU registers and memory,
 - viruses were not floating-point code
 - Viruses then began to deliberately use the floating point coprocessor registers and instructions
 - More recently, MMX and SSE vector graphics instructions are appearing in viruses
 - Modern emulators responded by emulating all registers and instructions

Also, can use undocumented x86 instructions to similar effect!

Anti-Emulation: Exceptions

The emulation environment is not always able to predict whether the next instruction will cause an exception (why?)

Viruses have been written to exploit this problem by putting part of their code in exception handlers, and then causing very subtle exceptions to occur

If part of the virus decompressor or decryptor code is in an exception handler, the emulator will fail to decrypt or decompress and emulation will try to execute garbage bytes

Anti-Emulation: Time and Logic Bombs

Arrange for virus to execute only...

- > at certain times of day (time bomb),
- or only under random conditions (logic bomb) that might be controlled by random number generation
- If the time/condition not met, the virus just transfers control to the infected host program and does no damage
- When executed in an emulator, the virus will probably be dormant and cannot be analyzed by the emulator

AV analysis forced to use a debugger

Anti-Emulation: Dynamic Code Length

- Code emulation is very slow (interpretation of an entire simulated machine environment)
 - Not a problem for most viruses, as most they start up pretty quickly and are written in tight code
- Anti-emulation Armoring: However, use huge loops of do-nothing instructions with a few real virus instructions thrown in & a high loop count
 - Emulation becomes too lengthy
 - Can also be done by encoding brute-force decryptors into the virus, which run billions of instructions before successfully decrypting it

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Anti-heuristic Viruses

- Scanning a PE file for virus code patterns is not always as simple as using regular expressions,
 - especially when detecting new viruses or new variants
 - E.g., metamorphic viruses
- Search heuristics have been developed that can find suspicious code without having exact patterns from a pattern database
 - Static heuristics are used in scanners to analyze executable files
 - Dynamic heuristics analyze code running under emulation
- In addition to anti-emulation techniques directed against the dynamic heuristics, virus writers learned to write viruses in such a way as to evade the static scanners

Anti-heuristic Techniques

- It is common for virus code to be appended, or placed in the last section in the PE file
 - Scanners are programmed to flag PE files in which the entry point is directed to the last section of the file
 - This can cause a false positive for self-extracting archives, in which the extractor code is often at the end, so other heuristics must be used in combination with this one
- Virus writers responded by adding some of their own data sections after their code section, so the entry point was not the last section any more,
 - e.g., the Resure virus
 - ▶Peter Szor, "Attacks on Win 32 Part 2" for more details

Anti-heuristic Techniques

- Another way to append virus code without having the entry point be in the last section of the PE file
 - place the beginning of the virus in the slack area at the end of the host program code section,
 - ...with a jump to the appended virus code at the end
- Emulators have been designed to detect jumping from one PE file section to another

Anti-heuristic Techniques

- The EPO (entry-point obscuring) techniques that we studied previously are also anti-heuristic techniques, as they make it hard for a scanner to detect that control passes to a virus
 - Import Address Table (IAT) replacement
 - Call hijacking
 - Replacing an arbitrary call with a jump into the virus
- Modern scanners cannot just look at entry points and the top and bottom of PE files
 - Scanning is getting more expensive as a result

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Anti-goat Viruses

- Anti-virus software often uses goat files, a.k.a sacrificial goats, which are dummy files whose infection will signal the presence of a virus
 - Short, simple files, of known content
 - Easy to find a virus within them
 - Scattered around the disk in various file types that are prone to infection, e.g. *.exe, *.vbs, *.com
 - Also in various sizes, as viruses often infect only files with a certain minimum size
- Detecting which goat files are infected, and which are not, helps identify the virus

Anti-goat Techniques

- An anti-goat virus will try to detect goat files and avoid infecting them
- Files are examined for goat file characteristics:
 - Lots of no-ops and do-nothing instructions
 - Clusters of files with sequential numbers in their names, e.g. abcd0001.vbs, abcd0002.vbs, etc.
- As with all armored virus techniques, the point is not to be the mythical "undetectable virus", but just to slow down detection and analysis

Retroviruses

- In nature, a retrovirus replicates in a different manner than normal viruses and is thus partly immune to many antiviral drugs
 - The HIV retrovirus attacks the immune system
- Computer retroviruses directly attack anti-virus software in an effort to make themselves immune
- Szor, section 6.3: "A retrovirus is a computer virus that specifically tries to bypass or hinder the operation of an antivirus, personal firewall, or other security programs."

Vulnerability to Retroviruses

- Most of us log on to our personal computers with administrative capabilities
 - More convenient than having to change your login every time you perform an administrative task
 - Everyone was an administrator under DOS
- This gives a retrovirus the same administrative capabilities, meaning that it can kill anti-virus processes, remove anti-virus files, etc., just as you could
- Paves the way for other viruses to work freely
 - Might be the only function of a given retrovirus

Retroviruses: Direct Attack on Security Software

 With admin capabilities, a retrovirus can kill the processes that it recognizes as being AV or firewall software, behavior blockers, etc.

A behavior blocker is a background process, from the OS or from an AV package, that prevents certain suspicious behaviors, such as changing the interrupt chain or doing a disk write to an existing executable file

- Sometimes, settings in firewalls and AV monitors can be changed to bypass them without killing them (stealthier than killing)
- Can also delete AV or firewall files from disk

Retroviruses: Attack on Security Software Files

Many AV programs maintain an integrity checking database full of checksums and file sizes for various system files

Retroviruses can attack this database:

Delete the files (not too stealthy)

Modify the files so that checksums and sizes must be recomputed for infected files; this hides the infection by storing the new sizes and checksums

 Replace the database with a modified database that prevents virus detection, causes virus misidentification, and even launches viruses (i.e. database is now a Trojan horse)

Example: Altering Integrity Database Entries

The IDEA.6155 virus was designed to infect *.COM and *.EXE files while escaping detection in the integrity database:

The integrity database checksum record "thisfile.exe 23f7e65b" would be altered to "lhisfile.exe 23f7e65b" after infection of thisfile.exe

 \blacktriangleright notice the change to the first character from "t" to "l"

AV integrity checker concludes, on its next scan, that thisfile.exe must be a new file, so it computes a checksum of the infected file and adds a new record to the integrity database: "thisfile.exe 269b7fc2"

Infected file now seems to have integrity

> Full scan will be required, searching for virus patterns, etc., to find the problem; full scans are not done often because of time constraints

Retroviruses: Indirect Attacks on AV Programs

- Older versions of AV programs sometimes placed an integrity check record at the end of a validated file, with encrypted checksums
 - Files with a mark did not need to be scanned again
 - Cut down on scanning time by only scanning modified or newly created files
- The Tequila virus removed this record from files it infected, so that infection would not cause an integrity check failure
- Simpler non-cryptographic checksums, such as CRC, can be defeated by appending a few bytes to an infected file that cause its new CRC checksum to match the old one;
 <u>Hybris</u> worm used this technique on PE files, which had to be restored (could not be repaired)

Retroviruses: Deterring AV Use

- A retrovirus can attack analysis tools used by anti-virus researchers
- A retrovirus can remain dormant until it detects AV software, then start damaging the system
- What would you do if the following message appeared on your screen:
 - WARNING! Infected System!
 - Virus will do no harm unless you install
 - anti-virus software. Then it will destroy
 - your files!

Assignment

 Read Szor, Chapter 7 through section 7.5 (virus encryption techniques) before the next lecture